

# The WebSocket Protocol

## Abstract

The WebSocket Protocol enables two-way communication between a client running untrusted code in a controlled environment to a remote host that has opted-in to communications from that code. The security model used for this is the origin-based security model commonly used by web browsers. The protocol consists of an opening handshake followed by basic message framing, layered over TCP. The goal of this technology is to provide a mechanism for browser-based applications that need two-way communication with servers that does not rely on opening multiple HTTP connections (e.g., using XMLHttpRequest or <iframe>s and long polling).

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# 1. Introduction

## 1.1. Background

*This section is non-normative.*

Historically, creating web applications that need bidirectional communication between a client and a server (e.g., instant messaging and gaming applications) has required an abuse of HTTP to poll the server for updates while sending upstream notifications as distinct HTTP calls [RFC6202].

This results in a variety of problems:

- The server is forced to use a number of different underlying TCP connections for each client: one for sending information to the client and a new one for each incoming message.
- The wire protocol has a high overhead, with each client-to-server message having an HTTP header.
- The client-side script is forced to maintain a mapping from the outgoing connections to the incoming connection to track replies.

A simpler solution would be to use a single TCP connection for traffic in both directions. This is what the WebSocket Protocol provides. Combined with the WebSocket API [WSAPI], it provides an alternative to HTTP polling for two-way communication from a web page to a remote server.

The same technique can be used for a variety of web applications: games, stock tickers, multiuser applications with simultaneous editing, user interfaces exposing server-side services in real time, etc.

The WebSocket Protocol is designed to supersede existing bidirectional communication technologies that use HTTP as a transport layer to benefit from existing infrastructure (proxies, filtering, authentication). Such technologies were implemented as trade-offs between efficiency and reliability because HTTP was not initially meant to be used for bidirectional communication (see [RFC6202] for further discussion). The WebSocket Protocol attempts to address the goals of existing bidirectional HTTP technologies in the context of the existing HTTP infrastructure; as such, it is designed to work over HTTP ports 80 and 443 as well as to support HTTP proxies and intermediaries, even if this implies some complexity specific to the current environment. However, the design does not limit WebSocket to HTTP, and future implementations could use a simpler handshake over a dedicated port without reinventing the entire protocol. This last point is important because the traffic patterns of interactive messaging do not closely match standard HTTP traffic and can induce unusual loads on some components.

## 1.2. Protocol Overview

*This section is non-normative.*

The protocol has two parts: a handshake and the data transfer.

The handshake from the client looks as follows:

```
GET /chat HTTP/1.1
Host: server.example.com
Upgrade: websocket
Connection: Upgrade
Sec-WebSocket-Key: dGhlIHNhbXBsZSBub25jZQ==
Origin: http://example.com
Sec-WebSocket-Protocol: chat, superchat
Sec-WebSocket-Version: 13
```

The handshake from the server looks as follows:

```
HTTP/1.1 101 Switching Protocols
Upgrade: websocket
Connection: Upgrade
Sec-WebSocket-Accept: s3pPLMBiTxaQ9kYGzzhZRbK+xOo=
Sec-WebSocket-Protocol: chat
```

The leading line from the client follows the Request-Line format. The leading line from the server follows the Status-Line format. The Request-Line and Status-Line productions are defined in [RFC2616].

An unordered set of header fields comes after the leading line in both cases. The meaning of these header fields is specified in Section 4 of this document. Additional header fields may also be present, such as cookies [RFC6265]. The format and parsing of headers is as defined in [RFC2616].

Once the client and server have both sent their handshakes, and if the handshake was successful, then the data transfer part starts. This is a two-way communication channel where each side can, independently from the other, send data at will.

After a successful handshake, clients and servers transfer data back and forth in conceptual units referred to in this specification as "messages". On the wire, a message is composed of one or more frames. The WebSocket message does not necessarily correspond to a particular network layer framing, as a fragmented message may be coalesced or split by an intermediary.

A frame has an associated type. Each frame belonging to the same message contains the same type of data. Broadly speaking, there are types for textual data (which is interpreted as UTF-8 [RFC3629] text), binary data (whose interpretation is left up to the application), and control frames (which are not intended to carry data for the application but instead for protocol-level signaling, such as to signal that the connection should be closed). This version of the protocol defines six frame types and leaves ten reserved for future use.

### 1.3. Opening Handshake

*This section is non-normative.*

The opening handshake is intended to be compatible with HTTP-based server-side software and intermediaries, so that a single port can be used by both HTTP clients talking to that server and WebSocket clients talking to that server. To this end, the WebSocket client's handshake is an HTTP Upgrade request:

```
GET /chat HTTP/1.1
Host: server.example.com
Upgrade: websocket
Connection: Upgrade
Sec-WebSocket-Key: dGhlIHNhbXBsZSBub25jZQ==
Origin: http://example.com
Sec-WebSocket-Protocol: chat, superchat
Sec-WebSocket-Version: 13
```

In compliance with [RFC2616], header fields in the handshake may be sent by the client in any order, so the order in which different header fields are received is not significant.

The "Request-URI" of the GET method [RFC2616] is used to identify the endpoint of the WebSocket connection, both to allow multiple domains to be served from one IP address and to allow multiple WebSocket endpoints to be served by a single server.

The client includes the hostname in the |Host| header field of its handshake as per [RFC2616], so that both the client and the server can verify that they agree on which host is in use.

Additional header fields are used to select options in the WebSocket Protocol. Typical options available in this version are the subprotocol selector (`|Sec-WebSocket-Protocol|`), list of extensions support by the client (`|Sec-WebSocket-Extensions|`), `|Origin|` header field, etc. The `|Sec-WebSocket-Protocol|` request-header field can be used to indicate what subprotocols (application-level protocols layered over the WebSocket Protocol) are acceptable to the client. The server selects one or none of the acceptable protocols and echoes that value in its handshake to indicate that it has selected that protocol.

```
Sec-WebSocket-Protocol: chat
```

The `|Origin|` header field [RFC6454] is used to protect against unauthorized cross-origin use of a WebSocket server by scripts using the WebSocket API in a web browser. The server is informed of the script origin generating the WebSocket connection request. If the server does not wish to accept connections from this origin, it can choose to reject the connection by sending an appropriate HTTP error code. This header field is sent by browser clients; for non-browser clients, this header field may be sent if it makes sense in the context of those clients.

Finally, the server has to prove to the client that it received the client's WebSocket handshake, so that the server doesn't accept connections that are not WebSocket connections. This prevents an attacker from tricking a WebSocket server by sending it carefully crafted packets using XMLHttpRequest [XMLHttpRequest] or a form submission.

To prove that the handshake was received, the server has to take two pieces of information and combine them to form a response. The first piece of information comes from the `|Sec-WebSocket-Key|` header field in the client handshake:

```
Sec-WebSocket-Key: dGhlIHNhbXBsZSBub25jZQ==
```

For this header field, the server has to take the value (as present in the header field, e.g., the base64-encoded [RFC4648] version minus any leading and trailing whitespace) and concatenate this with the Globally Unique Identifier (GUID, [RFC4122]) "258EAF5-E914-47DA-95CA-C5AB0DC85B11" in string form, which is unlikely to be used by network endpoints that do not understand the WebSocket Protocol. A SHA-1 hash (160 bits) [FIPS.180-3], base64-encoded (see Section 4 of [RFC4648]), of this concatenation is then returned in the server's handshake.

Concretely, if as in the example above, the `|Sec-WebSocket-Key|` header field had the value "dGhlIHNhbXBsZSBub25jZQ==", the server would concatenate the string "258EAF5-E914-47DA-95CA-C5AB0DC85B11" to form the string "dGhlIHNhbXBsZSBub25jZQ==258EAF5-E914-47DA-95CA-C5AB0DC85B11". The server would then take the SHA-1 hash of this, giving the value 0xb3 0x7a 0x4f 0x2c 0xc0 0x62 0x4f 0x16 0x90 0xf6 0x46 0x06 0xcf 0x38 0x59 0x45 0xb2 0xbe 0xc4 0xea. This value is then base64-encoded (see Section 4 of [RFC4648]), to give the value "s3pPLMBiTxaQ9kYGzzhZRbK+xOo=". This value would then be echoed in the `|Sec-WebSocket-Accept|` header field.

The handshake from the server is much simpler than the client handshake. The first line is an HTTP Status-Line, with the status code 101:

```
HTTP/1.1 101 Switching Protocols
```

Any status code other than 101 indicates that the WebSocket handshake has not completed and that the semantics of HTTP still apply. The headers follow the status code.

The `|Connection|` and `|Upgrade|` header fields complete the HTTP Upgrade. The `|Sec-WebSocket-Accept|` header field indicates whether the server is willing to accept the connection. If present, this header field must include a hash of the client's nonce sent in `|Sec-WebSocket-Key|` along with a predefined GUID. Any other value must not be interpreted as an acceptance of the connection by the server.

```

HTTP/1.1 101 Switching Protocols
Upgrade: websocket
Connection: Upgrade
Sec-WebSocket-Accept: s3pPLMBiTxaQ9kYGzzhZRbK+xOo=

```

These fields are checked by the WebSocket client for scripted pages. If the `|Sec-WebSocket-Accept|` value does not match the expected value, if the header field is missing, or if the HTTP status code is not 101, the connection will not be established, and WebSocket frames will not be sent.

Option fields can also be included. In this version of the protocol, the main option field is `|Sec-WebSocket-Protocol|`, which indicates the subprotocol that the server has selected. WebSocket clients verify that the server included one of the values that was specified in the WebSocket client's handshake. A server that speaks multiple subprotocols has to make sure it selects one based on the client's handshake and specifies it in its handshake.

```
Sec-WebSocket-Protocol: chat
```

The server can also set cookie-related option fields to *set* cookies, as described in [\[RFC6265\]](#).

## 1.4. Closing Handshake

*This section is non-normative.*

The closing handshake is far simpler than the opening handshake.

Either peer can send a control frame with data containing a specified control sequence to begin the closing handshake (detailed in [Section 5.5.1](#)). Upon receiving such a frame, the other peer sends a Close frame in response, if it hasn't already sent one. Upon receiving *that* control frame, the first peer then closes the connection, safe in the knowledge that no further data is forthcoming.

After sending a control frame indicating the connection should be closed, a peer does not send any further data; after receiving a control frame indicating the connection should be closed, a peer discards any further data received.

It is safe for both peers to initiate this handshake simultaneously.

The closing handshake is intended to complement the TCP closing handshake (FIN/ACK), on the basis that the TCP closing handshake is not always reliable end-to-end, especially in the presence of intercepting proxies and other intermediaries.

By sending a Close frame and waiting for a Close frame in response, certain cases are avoided where data may be unnecessarily lost. For instance, on some platforms, if a socket is closed with data in the receive queue, a RST packet is sent, which will then cause `recv()` to fail for the party that received the RST, even if there was data waiting to be read.

## 1.5. Design Philosophy

*This section is non-normative.*

The WebSocket Protocol is designed on the principle that there should be minimal framing (the only framing that exists is to make the protocol frame-based instead of stream-based and to support a distinction between Unicode text and binary frames). It is expected that metadata would be layered on top of WebSocket by the application layer, in the same way that metadata is layered on top of TCP by the application layer (e.g., HTTP).

Conceptually, WebSocket is really just a layer on top of TCP that does the following:

- adds a web origin-based security model for browsers
- adds an addressing and protocol naming mechanism to support multiple services on one port and multiple host names on one IP address



- layers a framing mechanism on top of TCP to get back to the IP packet mechanism that TCP is built on, but without length limits
- includes an additional closing handshake in-band that is designed to work in the presence of proxies and other intermediaries

Other than that, WebSocket adds nothing. Basically it is intended to be as close to just exposing raw TCP to script as possible given the constraints of the Web. It's also designed in such a way that its servers can share a port with HTTP servers, by having its handshake be a valid HTTP Upgrade request. One could conceptually use other protocols to establish client-server messaging, but the intent of WebSockets is to provide a relatively simple protocol that can coexist with HTTP and deployed HTTP infrastructure (such as proxies) and that is as close to TCP as is safe for use with such infrastructure given security considerations, with targeted additions to simplify usage and keep simple things simple (such as the addition of message semantics).

The protocol is intended to be extensible; future versions will likely introduce additional concepts such as multiplexing.

## 1.6. Security Model

*This section is non-normative.*

The WebSocket Protocol uses the origin model used by web browsers to restrict which web pages can contact a WebSocket server when the WebSocket Protocol is used from a web page. Naturally, when the WebSocket Protocol is used by a dedicated client directly (i.e., not from a web page through a web browser), the origin model is not useful, as the client can provide any arbitrary origin string.

This protocol is intended to fail to establish a connection with servers of pre-existing protocols like SMTP [RFC5321] and HTTP, while allowing HTTP servers to opt-in to supporting this protocol if desired. This is achieved by having a strict and elaborate handshake and by limiting the data that can be inserted into the connection before the handshake is finished (thus limiting how much the server can be influenced).

It is similarly intended to fail to establish a connection when data from other protocols, especially HTTP, is sent to a WebSocket server, for example, as might happen if an HTML "form" were submitted to a WebSocket server. This is primarily achieved by requiring that the server prove that it read the handshake, which it can only do if the handshake contains the appropriate parts, which can only be sent by a WebSocket client. In particular, at the time of writing of this specification, fields starting with |Sec-| cannot be set by an attacker from a web browser using only HTML and JavaScript APIs such as XMLHttpRequest [XMLHttpRequest].

## 1.7. Relationship to TCP and HTTP

*This section is non-normative.*

The WebSocket Protocol is an independent TCP-based protocol. Its only relationship to HTTP is that its handshake is interpreted by HTTP servers as an Upgrade request.

By default, the WebSocket Protocol uses port 80 for regular WebSocket connections and port 443 for WebSocket connections tunneled over Transport Layer Security (TLS) [RFC2818].

## 1.8. Establishing a Connection

*This section is non-normative.*

When a connection is to be made to a port that is shared by an HTTP server (a situation that is quite likely to occur with traffic to ports 80 and 443), the connection will appear to the HTTP server to be a regular GET request with an Upgrade offer. In relatively simple setups with just one IP address and a single server for all traffic to a single hostname, this might allow a practical way for systems based on the WebSocket Protocol to be deployed. In more elaborate setups (e.g., with load balancers and multiple servers), a dedicated set of hosts for WebSocket connections separate from the HTTP servers is probably easier to manage. At the time of writing of this specification, it should be noted that connections on ports 80 and 443 have significantly different

success rates, with connections on port 443 being significantly more likely to succeed, though this may change with time.

## 1.9. Subprotocols Using the WebSocket Protocol

*This section is non-normative.*

The client can request that the server use a specific subprotocol by including the |Sec-WebSocket-Protocol| field in its handshake. If it is specified, the server needs to include the same field and one of the selected subprotocol values in its response for the connection to be established.

These subprotocol names should be registered as per [Section 11.5](#). To avoid potential collisions, it is recommended to use names that contain the ASCII version of the domain name of the subprotocol's originator. For example, if Example Corporation were to create a Chat subprotocol to be implemented by many servers around the Web, they could name it "chat.example.com". If the Example Organization called their competing subprotocol "chat.example.org", then the two subprotocols could be implemented by servers simultaneously, with the server dynamically selecting which subprotocol to use based on the value sent by the client.

Subprotocols can be versioned in backward-incompatible ways by changing the subprotocol name, e.g., going from "bookings.example.net" to "v2.bookings.example.net". These subprotocols would be considered completely separate by WebSocket clients. Backward-compatible versioning can be implemented by reusing the same subprotocol string but carefully designing the actual subprotocol to support this kind of extensibility.

## 2. Conformance Requirements

All diagrams, examples, and notes in this specification are non-normative, as are all sections explicitly marked non-normative. Everything else in this specification is normative.

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

Requirements phrased in the imperative as part of algorithms (such as "strip any leading space characters" or "return false and abort these steps") are to be interpreted with the meaning of the key word ("MUST", "SHOULD", "MAY", etc.) used in introducing the algorithm.

Conformance requirements phrased as algorithms or specific steps MAY be implemented in any manner, so long as the end result is equivalent. (In particular, the algorithms defined in this specification are intended to be easy to follow and not intended to be performant.)

### 2.1. Terminology and Other Conventions

*ASCII* shall mean the character-encoding scheme defined in [ANSI.X3-4.1986].

This document makes reference to UTF-8 values and uses UTF-8 notational formats as defined in STD 63 [RFC3629].

Key terms such as named algorithms or definitions are indicated like *this*.

Names of header fields or variables are indicated like |this|.

Variable values are indicated like /this/.

This document references the procedure to *Fail the WebSocket Connection*. This procedure is defined in Section 7.1.7.

*Converting a string to ASCII lowercase* means replacing all characters in the range U+0041 to U+005A (i.e., LATIN CAPITAL LETTER A to LATIN CAPITAL LETTER Z) with the corresponding characters in the range U+0061 to U+007A (i.e., LATIN SMALL LETTER A to LATIN SMALL LETTER Z).

Comparing two strings in an *ASCII case-insensitive* manner means comparing them exactly, code point for code point, except that the characters in the range U+0041 to U+005A (i.e., LATIN CAPITAL LETTER A to LATIN CAPITAL LETTER Z) and the corresponding characters in the range U+0061 to U+007A (i.e., LATIN SMALL LETTER A to LATIN SMALL LETTER Z) are considered to also match.

The term "URI" is used in this document as defined in [RFC3986].

When an implementation is required to *send* data as part of the WebSocket Protocol, the implementation MAY delay the actual transmission arbitrarily, e.g., buffering data so as to send fewer IP packets.

Note that this document uses both [RFC5234] and [RFC2616] variants of ABNF in different sections.

### 3. WebSocket URIs

This specification defines two URI schemes, using the ABNF syntax defined in RFC 5234 [RFC5234], and terminology and ABNF productions defined by the URI specification RFC 3986 [RFC3986].

```
ws-URI = "ws:" "://" host [ ":" port ] path [ "?" query ]
wss-URI = "wss:" "://" host [ ":" port ] path [ "?" query ]

host = <host, defined in [RFC3986], Section 3.2.2>
port = <port, defined in [RFC3986], Section 3.2.3>
path = <path-abempty, defined in [RFC3986], Section 3.3>
query = <query, defined in [RFC3986], Section 3.4>
```

The port component is OPTIONAL; the default for "ws" is port 80, while the default for "wss" is port 443.

The URI is called "secure" (and it is said that "the secure flag is set") if the scheme component matches "wss" case-insensitively.

The "resource-name" (also known as /resource name/ in [Section 4.1](#)) can be constructed by concatenating the following:

- "/" if the path component is empty
- the path component
- "?" if the query component is non-empty
- the query component

Fragment identifiers are meaningless in the context of WebSocket URIs and MUST NOT be used on these URIs. As with any URI scheme, the character "#", when not indicating the start of a fragment, MUST be escaped as %23.

## 4. Opening Handshake

### 4.1. Client Requirements

To *Establish a WebSocket Connection*, a client opens a connection and sends a handshake as defined in this section. A connection is defined to initially be in a `CONNECTING` state. A client will need to supply a `/host/`, `/port/`, `/resource name/`, and a `/secure/` flag, which are the components of a WebSocket URI as discussed in [Section 3](#), along with a list of `/protocols/` and `/extensions/` to be used. Additionally, if the client is a web browser, it supplies `/origin/`.

Clients running in controlled environments, e.g., browsers on mobile handsets tied to specific carriers, MAY offload the management of the connection to another agent on the network. In such a situation, the client for the purposes of this specification is considered to include both the handset software and any such agents.

When the client is to *Establish a WebSocket Connection* given a set of `(/host/, /port/, /resource name/, and /secure/ flag)`, along with a list of `/protocols/` and `/extensions/` to be used, and an `/origin/` in the case of web browsers, it **MUST** open a connection, send an opening handshake, and read the server's handshake in response. The exact requirements of how the connection should be opened, what should be sent in the opening handshake, and how the server's response should be interpreted are as follows in this section. In the following text, we will use terms from [Section 3](#), such as `"/host/"` and `"/secure/ flag"` as defined in that section.

1. The components of the WebSocket URI passed into this algorithm (`/host/`, `/port/`, `/resource name/`, and `/secure/ flag`) **MUST** be valid according to the specification of WebSocket URIs specified in [Section 3](#). If any of the components are invalid, the client **MUST** *Fail the WebSocket Connection* and abort these steps.
2. If the client already has a WebSocket connection to the remote host (IP address) identified by `/host/` and port `/port/` pair, even if the remote host is known by another name, the client **MUST** wait until that connection has been established or for that connection to have failed. There **MUST** be no more than one connection in a `CONNECTING` state. If multiple connections to the same IP address are attempted simultaneously, the client **MUST** serialize them so that there is no more than one connection at a time running through the following steps.

If the client cannot determine the IP address of the remote host (for example, because all communication is being done through a proxy server that performs DNS queries itself), then the client **MUST** assume for the purposes of this step that each host name refers to a distinct remote host, and instead the client **SHOULD** limit the total number of simultaneous pending connections to a reasonably low number (e.g., the client might allow simultaneous pending connections to `a.example.com` and `b.example.com`, but if thirty simultaneous connections to a single host are requested, that may not be allowed). For example, in a web browser context, the client needs to consider the number of tabs the user has open in setting a limit to the number of simultaneous pending connections.

NOTE: This makes it harder for a script to perform a denial-of-service attack by just opening a large number of WebSocket connections to a remote host. A server can further reduce the load on itself when attacked by pausing before closing the connection, as that will reduce the rate at which the client reconnects.

NOTE: There is no limit to the number of established WebSocket connections a client can have with a single remote host. Servers can refuse to accept connections from hosts/IP addresses with an excessive number of existing connections or disconnect resource-hogging connections when suffering high load.

3. *Proxy Usage*: If the client is configured to use a proxy when using the WebSocket Protocol to connect to host `/host/` and port `/port/`, then the client **SHOULD** connect to that proxy and ask it to open a TCP connection to the host given by `/host/` and the port given by `/port/`.

EXAMPLE: For example, if the client uses an HTTP proxy for all traffic, then if it was to try to connect to port 80 on server `example.com`, it might send the following lines to the proxy server:

```
CONNECT example.com:80 HTTP/1.1
Host: example.com
```

If there was a password, the connection might look like:

```
CONNECT example.com:80 HTTP/1.1
Host: example.com
Proxy-authorization: Basic ZWRuYW1vZGU6bm9jYXB1cyE=
```

If the client is not configured to use a proxy, then a direct TCP connection SHOULD be opened to the host given by /host/ and the port given by /port/.

NOTE: Implementations that do not expose explicit UI for selecting a proxy for WebSocket connections separate from other proxies are encouraged to use a SOCKS5 [RFC1928] proxy for WebSocket connections, if available, or failing that, to prefer the proxy configured for HTTPS connections over the proxy configured for HTTP connections.

For the purpose of proxy autoconfiguration scripts, the URI to pass the function MUST be constructed from /host/, /port/, /resource name/, and the /secure/ flag using the definition of a WebSocket URI as given in Section 3.

NOTE: The WebSocket Protocol can be identified in proxy autoconfiguration scripts from the scheme ("ws" for unencrypted connections and "wss" for encrypted connections).

4. If the connection could not be opened, either because a direct connection failed or because any proxy used returned an error, then the client MUST *Fail the WebSocket Connection* and abort the connection attempt.
5. If /secure/ is true, the client MUST perform a TLS handshake over the connection after opening the connection and before sending the handshake data [RFC2818]. If this fails (e.g., the server's certificate could not be verified), then the client MUST *Fail the WebSocket Connection* and abort the connection. Otherwise, all further communication on this channel MUST run through the encrypted tunnel [RFC5246]. Clients MUST use the Server Name Indication extension in the TLS handshake [RFC6066].

Once a connection to the server has been established (including a connection via a proxy or over a TLS-encrypted tunnel), the client MUST send an opening handshake to the server. The handshake consists of an HTTP Upgrade request, along with a list of required and optional header fields. The requirements for this handshake are as follows.

1. The handshake MUST be a valid HTTP request as specified by [RFC2616].
2. The method of the request MUST be GET, and the HTTP version MUST be at least 1.1.  
For example, if the WebSocket URI is "ws://example.com/chat", the first line sent should be "GET /chat HTTP/1.1".
3. The "Request-URI" part of the request MUST match the /resource name/ defined in Section 3 (a relative URI) or be an absolute http/https URI that, when parsed, has a /resource name/, /host/, and /port/ that match the corresponding ws/wss URI.
4. The request MUST contain a |Host| header field whose value contains /host/ plus optionally ":" followed by /port/ (when not using the default port).
5. The request MUST contain an |Upgrade| header field whose value MUST include the "websocket" keyword.
6. The request MUST contain a |Connection| header field whose value MUST include the "Upgrade" token.
7. The request MUST include a header field with the name |Sec-WebSocket-Key|. The value of this header field MUST be a nonce consisting of a randomly selected 16-byte value that has been base64-encoded (see Section 4 of [RFC4648]). The nonce MUST be selected randomly for each connection.  
NOTE: As an example, if the randomly selected value was the sequence of bytes 0x01 0x02 0x03 0x04 0x05 0x06 0x07 0x08 0x09 0x0a 0x0b 0x0c 0x0d 0x0e 0x0f 0x10, the value of the header field would be "AQIDBAUGBwgJCgsMDQ4PEc=="
8. The request MUST include a header field with the name |Origin| [RFC6454] if the request is coming from a browser client. If the connection is from a non-browser client, the request MAY include this header field if the semantics of that client match the use-case described here for browser clients. The value of this header

field is the ASCII serialization of origin of the context in which the code establishing the connection is running. See [RFC6454] for the details of how this header field value is constructed.

As an example, if code downloaded from www.example.com attempts to establish a connection to ww2.example.com, the value of the header field would be "http://www.example.com".

9. The request MUST include a header field with the name |Sec-WebSocket-Version|. The value of this header field MUST be 13.  
NOTE: Although draft versions of this document (-09, -10, -11, and -12) were posted (they were mostly comprised of editorial changes and clarifications and not changes to the wire protocol), values 9, 10, 11, and 12 were not used as valid values for Sec-WebSocket-Version. These values were reserved in the IANA registry but were not and will not be used.
10. The request MAY include a header field with the name |Sec-WebSocket-Protocol|. If present, this value indicates one or more comma-separated subprotocol the client wishes to speak, ordered by preference. The elements that comprise this value MUST be non-empty strings with characters in the range U+0021 to U+007E not including separator characters as defined in [RFC2616] and MUST all be unique strings. The ABNF for the value of this header field is 1#token, where the definitions of constructs and rules are as given in [RFC2616].
11. The request MAY include a header field with the name |Sec-WebSocket-Extensions|. If present, this value indicates the protocol-level extension(s) the client wishes to speak. The interpretation and format of this header field is described in Section 9.1.
12. The request MAY include any other header fields, for example, cookies [RFC6265] and/or authentication-related header fields such as the |Authorization| header field [RFC2616], which are processed according to documents that define them.

Once the client's opening handshake has been sent, the client MUST wait for a response from the server before sending any further data. The client MUST validate the server's response as follows:

1. If the status code received from the server is not 101, the client handles the response per HTTP [RFC2616] procedures. In particular, the client might perform authentication if it receives a 401 status code; the server might redirect the client using a 3xx status code (but clients are not required to follow them), etc. Otherwise, proceed as follows.
2. If the response lacks an |Upgrade| header field or the |Upgrade| header field contains a value that is not an ASCII case-insensitive match for the value "websocket", the client MUST *Fail the WebSocket Connection*.
3. If the response lacks a |Connection| header field or the |Connection| header field doesn't contain a token that is an ASCII case-insensitive match for the value "Upgrade", the client MUST *Fail the WebSocket Connection*.
4. If the response lacks a |Sec-WebSocket-Accept| header field or the |Sec-WebSocket-Accept| contains a value other than the base64-encoded SHA-1 of the concatenation of the |Sec-WebSocket-Key| (as a string, not base64-decoded) with the string "258EAF5-E914-47DA-95CA-C5AB0DC85B11" but ignoring any leading and trailing whitespace, the client MUST *Fail the WebSocket Connection*.
5. If the response includes a |Sec-WebSocket-Extensions| header field and this header field indicates the use of an extension that was not present in the client's handshake (the server has indicated an extension not requested by the client), the client MUST *Fail the WebSocket Connection*. (The parsing of this header field to determine which extensions are requested is discussed in Section 9.1.)
6. If the response includes a |Sec-WebSocket-Protocol| header field and this header field indicates the use of a subprotocol that was not present in the client's handshake (the server has indicated a subprotocol not requested by the client), the client MUST *Fail the WebSocket Connection*.

If the server's response does not conform to the requirements for the server's handshake as defined in this section and in Section 4.2.2, the client MUST *Fail the WebSocket Connection*.

Please note that according to [RFC2616], all header field names in both HTTP requests and HTTP responses are case-insensitive.

If the server's response is validated as provided for above, it is said that *The WebSocket Connection is Established* and that the WebSocket Connection is in the OPEN state. The *Extensions In Use* is defined to

be a (possibly empty) string, the value of which is equal to the value of the `|Sec-WebSocket-Extensions|` header field supplied by the server's handshake or the null value if that header field was not present in the server's handshake. The *Subprotocol In Use* is defined to be the value of the `|Sec-WebSocket-Protocol|` header field in the server's handshake or the null value if that header field was not present in the server's handshake. Additionally, if any header fields in the server's handshake indicate that cookies should be set (as defined by [\[RFC6265\]](#)), these cookies are referred to as *Cookies Set During the Server's Opening Handshake*.

## 4.2. Server-Side Requirements

Servers MAY offload the management of the connection to other agents on the network, for example, load balancers and reverse proxies. In such a situation, the server for the purposes of this specification is considered to include all parts of the server-side infrastructure from the first device to terminate the TCP connection all the way to the server that processes requests and sends responses.

EXAMPLE: A data center might have a server that responds to WebSocket requests with an appropriate handshake and then passes the connection to another server to actually process the data frames. For the purposes of this specification, the "server" is the combination of both computers.

### 4.2.1. Reading the Client's Opening Handshake

When a client starts a WebSocket connection, it sends its part of the opening handshake. The server must parse at least part of this handshake in order to obtain the necessary information to generate the server part of the handshake.

The client's opening handshake consists of the following parts. If the server, while reading the handshake, finds that the client did not send a handshake that matches the description below (note that as per [\[RFC2616\]](#), the order of the header fields is not important), including but not limited to any violations of the ABNF grammar specified for the components of the handshake, the server MUST stop processing the client's handshake and return an HTTP response with an appropriate error code (such as 400 Bad Request).

1. An HTTP/1.1 or higher GET request, including a "Request-URI" [\[RFC2616\]](#) that should be interpreted as a `/resource name/` defined in [Section 3](#) (or an absolute HTTP/HTTPS URI containing the `/resource name/`).
2. A `|Host|` header field containing the server's authority.
3. An `|Upgrade|` header field containing the value "websocket", treated as an ASCII case-insensitive value.
4. A `|Connection|` header field that includes the token "Upgrade", treated as an ASCII case-insensitive value.
5. A `|Sec-WebSocket-Key|` header field with a base64-encoded (see [Section 4](#) of [\[RFC4648\]](#)) value that, when decoded, is 16 bytes in length.
6. A `|Sec-WebSocket-Version|` header field, with a value of 13.
7. Optionally, an `|Origin|` header field. This header field is sent by all browser clients. A connection attempt lacking this header field SHOULD NOT be interpreted as coming from a browser client.
8. Optionally, a `|Sec-WebSocket-Protocol|` header field, with a list of values indicating which protocols the client would like to speak, ordered by preference.
9. Optionally, a `|Sec-WebSocket-Extensions|` header field, with a list of values indicating which extensions the client would like to speak. The interpretation of this header field is discussed in [Section 9.1](#).
10. Optionally, other header fields, such as those used to send cookies or request authentication to a server. Unknown header fields are ignored, as per [\[RFC2616\]](#).

### 4.2.2. Sending the Server's Opening Handshake

When a client establishes a WebSocket connection to a server, the server MUST complete the following steps to accept the connection and send the server's opening handshake.

1. If the connection is happening on an HTTPS (HTTP-over-TLS) port, perform a TLS handshake over the connection. If this fails (e.g., the client indicated a host name in the extended client hello "server\_name" extension that the server does not host), then close the connection; otherwise, all further communication for the connection (including the server's handshake) MUST run through the encrypted tunnel [\[RFC5246\]](#).



2. The server can perform additional client authentication, for example, by returning a 401 status code with the corresponding |WWW-Authenticate| header field as described in [RFC2616].
3. The server MAY redirect the client using a 3xx status code [RFC2616]. Note that this step can happen together with, before, or after the optional authentication step described above.
4. Establish the following information:

/origin/	The  Origin  header field in the client's handshake indicates the origin of the script establishing the connection. The origin is serialized to ASCII and converted to lowercase. The server MAY use this information as part of a determination of whether to accept the incoming connection. If the server does not validate the origin, it will accept connections from anywhere. If the server does not wish to accept this connection, it MUST return an appropriate HTTP error code (e.g., 403 Forbidden) and abort the WebSocket handshake described in this section. For more detail, refer to <a href="#">Section 10</a> .
/key/	The  Sec-WebSocket-Key  header field in the client's handshake includes a base64-encoded value that, if decoded, is 16 bytes in length. This (encoded) value is used in the creation of the server's handshake to indicate an acceptance of the connection. It is not necessary for the server to base64-decode the  Sec-WebSocket-Key  value.
/version/	The  Sec-WebSocket-Version  header field in the client's handshake includes the version of the WebSocket Protocol with which the client is attempting to communicate. If this version does not match a version understood by the server, the server MUST abort the WebSocket handshake described in this section and instead send an appropriate HTTP error code (such as 426 Upgrade Required) and a  Sec-WebSocket-Version  header field indicating the version(s) the server is capable of understanding.
/resource name/	An identifier for the service provided by the server. If the server provides multiple services, then the value should be derived from the resource name given in the client's handshake in the "Request-URI" [RFC2616] of the GET method. If the requested service is not available, the server MUST send an appropriate HTTP error code (such as 404 Not Found) and abort the WebSocket handshake.
/subprotocol/	Either a single value representing the subprotocol the server is ready to use or null. The value chosen MUST be derived from the client's handshake, specifically by selecting one of the values from the  Sec-WebSocket-Protocol  field that the server is willing to use for this connection (if any). If the client's handshake did not contain such a header field or if the server does not agree to any of the client's requested subprotocols, the only acceptable value is null. The absence of such a field is equivalent to the null value (meaning that if the server does not wish to agree to one of the suggested subprotocols, it MUST NOT send back a  Sec-WebSocket-Protocol  header field in its response). The empty string is not the same as the null value for these purposes and is not a legal value for this field. The ABNF for the value of this header field is (token), where the definitions of constructs and rules are as given in [RFC2616].
/extensions/	A (possibly empty) list representing the protocol-level extensions the server is ready to use. If the server supports multiple extensions, then the value MUST be derived from the client's handshake, specifically by selecting one or more of the values from the  Sec-WebSocket-Extensions  field. The absence of such a field is equivalent to the null value. The empty string is not the same as the null value for these purposes. Extensions not listed by the client MUST

NOT be listed. The method by which these values should be selected and interpreted is discussed in [Section 9.1](#).

5. If the server chooses to accept the incoming connection, it MUST reply with a valid HTTP response indicating the following.
  1. A Status-Line with a 101 response code as per RFC 2616 [RFC2616]. Such a response could look like "HTTP/1.1 101 Switching Protocols".
  2. An |Upgrade| header field with value "websocket" as per RFC 2616 [RFC2616].
  3. A |Connection| header field with value "Upgrade".
  4. A |Sec-WebSocket-Accept| header field. The value of this header field is constructed by concatenating /key/, defined above in step 4 in [Section 4.2.2](#), with the string "258EAF5E914-47DA-95CA-C5AB0DC85B11", taking the SHA-1 hash of this concatenated value to obtain a 20-byte value and base64-encoding (see Section 4 of [RFC4648]) this 20-byte hash. The ABNF [RFC2616] of this header field is defined as follows:

```

Sec-WebSocket-Accept      = base64-value-non-empty
base64-value-non-empty  = (1*base64-data [ base64-padding ] ) |
                           base64-padding
base64-data              = 4base64-character
base64-padding          = (2base64-character "==") |
                           (3base64-character "=")
base64-character        = ALPHA | DIGIT | "+" | "/"

```

NOTE: As an example, if the value of the |Sec-WebSocket-Key| header field in the client's handshake were "dGhIHNhbXBsZSBub25jZQ==", the server would append the string "258EAF5E914-47DA-95CA-C5AB0DC85B11" to form the string "dGhIHNhbXBsZSBub25jZQ==258EAF5E914-47DA-95CA-C5AB0DC85B11". The server would then take the SHA-1 hash of this string, giving the value 0xb3 0x7a 0x4f 0x2c 0xc0 0x62 0x4f 0x16 0x90 0xf6 0x46 0x06 0xcf 0x38 0x59 0x45 0xb2 0xbe 0xc4 0xea. This value is then base64-encoded, to give the value "s3pPLMBiTxaQ9kYGzZhZRbK+xOo=", which would be returned in the |Sec-WebSocket-Accept| header field.

5. Optionally, a |Sec-WebSocket-Protocol| header field, with a value /subprotocol/ as defined in step 4 in [Section 4.2.2](#).
6. Optionally, a |Sec-WebSocket-Extensions| header field, with a value /extensions/ as defined in step 4 in [Section 4.2.2](#). If multiple extensions are to be used, they can all be listed in a single |Sec-WebSocket-Extensions| header field or split between multiple instances of the |Sec-WebSocket-Extensions| header field.

This completes the server's handshake. If the server finishes these steps without aborting the WebSocket handshake, the server considers the WebSocket connection to be established and that the WebSocket connection is in the OPEN state. At this point, the server may begin sending (and receiving) data.

### 4.3. Collected ABNF for New Header Fields Used in Handshake

This section is using ABNF syntax/rules from Section 2.1 of [RFC2616], including the "implied \*LWS rule".

Note that the following ABNF conventions are used in this section. Some names of the rules correspond to names of the corresponding header fields. Such rules express values of the corresponding header fields, for example, the Sec-WebSocket-Key ABNF rule describes syntax of the |Sec-WebSocket-Key| header field value. ABNF rules with the "-Client" suffix in the name are only used in requests sent by the client to the server; ABNF rules with the "-Server" suffix in the name are only used in responses sent by the server to the client. For example, the ABNF rule Sec-WebSocket-Protocol-Client describes syntax of the |Sec-WebSocket-Protocol| header field value sent by the client to the server.

The following new header fields can be sent during the handshake from the client to the server:

```

Sec-WebSocket-Key = base64-value-non-empty
Sec-WebSocket-Extensions = extension-list
Sec-WebSocket-Protocol-Client = 1#token
Sec-WebSocket-Version-Client = version

base64-value-non-empty = (1*base64-data [ base64-padding ] ) |
                          base64-padding
base64-data             = 4base64-character
base64-padding         = (2base64-character "==") |
                          (3base64-character "=")
base64-character       = ALPHA | DIGIT | "+" | "/"
extension-list         = 1#extension
extension               = extension-token *( ";" extension-param )
extension-token        = registered-token
registered-token       = token
extension-param        = token [ "=" (token | quoted-string) ]
                        ; When using the quoted-string syntax variant, the value
                        ; after quoted-string unescaping MUST conform to the
                        ; 'token' ABNF.
NZDIGIT                = "1" | "2" | "3" | "4" | "5" | "6" |
                          "7" | "8" | "9"
version                 = DIGIT | (NZDIGIT DIGIT) |
                          ("1" DIGIT DIGIT) | ("2" DIGIT DIGIT)
                        ; Limited to 0-255 range, with no leading zeros

```

The following new header fields can be sent during the handshake from the server to the client:

```

Sec-WebSocket-Extensions = extension-list
Sec-WebSocket-Accept      = base64-value-non-empty
Sec-WebSocket-Protocol-Server = token
Sec-WebSocket-Version-Server = 1#version

```

#### 4.4. Supporting Multiple Versions of WebSocket Protocol

This section provides some guidance on supporting multiple versions of the WebSocket Protocol in clients and servers.

Using the WebSocket version advertisement capability (the |Sec-WebSocket-Version| header field), a client can initially request the version of the WebSocket Protocol that it prefers (which doesn't necessarily have to be the latest supported by the client). If the server supports the requested version and the handshake message is otherwise valid, the server will accept that version. If the server doesn't support the requested version, it **MUST** respond with a |Sec-WebSocket-Version| header field (or multiple |Sec-WebSocket-Version| header fields) containing all versions it is willing to use. At this point, if the client supports one of the advertised versions, it can repeat the WebSocket handshake using a new version value.

The following example demonstrates version negotiation described above:

```
GET /chat HTTP/1.1
Host: server.example.com
Upgrade: websocket
Connection: Upgrade
...
Sec-WebSocket-Version: 25
```

The response from the server might look as follows:

```
HTTP/1.1 400 Bad Request
...
Sec-WebSocket-Version: 13, 8, 7
```

Note that the last response from the server might also look like:

```
HTTP/1.1 400 Bad Request
...
Sec-WebSocket-Version: 13
Sec-WebSocket-Version: 8, 7
```

The client now repeats the handshake that conforms to version 13:

```
GET /chat HTTP/1.1
Host: server.example.com
Upgrade: websocket
Connection: Upgrade
...
Sec-WebSocket-Version: 13
```



Opcode:	4 bits Defines the interpretation of the "Payload data". If an unknown opcode is received, the receiving endpoint <i>MUST Fail the WebSocket Connection</i> . The following values are defined. <ul style="list-style-type: none"> <li>• %x0 denotes a continuation frame</li> <li>• %x1 denotes a text frame</li> <li>• %x2 denotes a binary frame</li> <li>• %x3-7 are reserved for further non-control frames</li> <li>• %x8 denotes a connection close</li> <li>• %x9 denotes a ping</li> <li>• %xA denotes a pong</li> <li>• %xB-F are reserved for further control frames</li> </ul>
Mask:	1 bit Defines whether the "Payload data" is masked. If set to 1, a masking key is present in masking-key, and this is used to unmask the "Payload data" as per <a href="#">Section 5.3</a> . All frames sent from client to server have this bit set to 1.
Payload length:	7 bits, 7+16 bits, or 7+64 bits The length of the "Payload data", in bytes: if 0-125, that is the payload length. If 126, the following 2 bytes interpreted as a 16-bit unsigned integer are the payload length. If 127, the following 8 bytes interpreted as a 64-bit unsigned integer (the most significant bit <i>MUST</i> be 0) are the payload length. Multibyte length quantities are expressed in network byte order. Note that in all cases, the minimal number of bytes <i>MUST</i> be used to encode the length, for example, the length of a 124-byte-long string can't be encoded as the sequence 126, 0, 124. The payload length is the length of the "Extension data" + the length of the "Application data". The length of the "Extension data" may be zero, in which case the payload length is the length of the "Application data".
Masking-key:	0 or 4 bytes All frames sent from the client to the server are masked by a 32#bit value that is contained within the frame. This field is present if the mask bit is set to 1 and is absent if the mask bit is set to 0. See <a href="#">Section 5.3</a> for further information on client-to-server masking.
Payload data:	(x+y) bytes The "Payload data" is defined as "Extension data" concatenated with "Application data".
Extension data:	x bytes The "Extension data" is 0 bytes unless an extension has been negotiated. Any extension <i>MUST</i> specify the length of the "Extension data", or how that length may be calculated, and how the extension use <i>MUST</i> be negotiated during the opening handshake. If present, the "Extension data" is included in the total payload length.
Application data:	y bytes Arbitrary "Application data", taking up the remainder of the frame after any "Extension data". The length of the "Application data" is equal to the payload length minus the length of the "Extension data".

The base framing protocol is formally defined by the following ABNF [\[RFC5234\]](#). It is important to note that the representation of this data is binary, not ASCII characters. As such, a field with a length of 1 bit that takes values %x0 / %x1 is represented as a single bit whose value is 0 or 1, not a full byte (octet) that stands for the

characters "0" or "1" in the ASCII encoding. A field with a length of 4 bits with values between %x0-F again is represented by 4 bits, again NOT by an ASCII character or full byte (octet) with these values. [RFC5234] does not specify a character encoding: "Rules resolve into a string of terminal values, sometimes called characters. In ABNF, a character is merely a non-negative integer. In certain contexts, a specific mapping (encoding) of values into a character set (such as ASCII) will be specified." Here, the specified encoding is a binary encoding where each terminal value is encoded in the specified number of bits, which varies for each field.

```

ws-frame          = frame-fin          ; 1 bit in length
                  frame-rsv1          ; 1 bit in length
                  frame-rsv2          ; 1 bit in length
                  frame-rsv3          ; 1 bit in length
                  frame-opcode        ; 4 bits in length
                  frame-masked        ; 1 bit in length
                  frame-payload-length ; either 7, 7+16,
                                      ; or 7+64 bits in
                                      ; length
                  [ frame-masking-key ] ; 32 bits in length
                  frame-payload-data  ; n*8 bits in
                                      ; length, where
                                      ; n >= 0

frame-fin         = %x0 ; more frames of this message follow
                  / %x1 ; final frame of this message
                  ; 1 bit in length

frame-rsv1        = %x0 / %x1
                  ; 1 bit in length, MUST be 0 unless
                  ; negotiated otherwise

frame-rsv2        = %x0 / %x1
                  ; 1 bit in length, MUST be 0 unless
                  ; negotiated otherwise

frame-rsv3        = %x0 / %x1
                  ; 1 bit in length, MUST be 0 unless
                  ; negotiated otherwise

frame-opcode      = frame-opcode-non-control /
                  frame-opcode-control /
                  frame-opcode-cont

frame-opcode-cont = %x0 ; frame continuation

frame-opcode-non-control = %x1 ; text frame
                          / %x2 ; binary frame
                          / %x3-7
                          ; 4 bits in length,
                          ; reserved for further non-control frames

frame-opcode-control = %x8 ; connection close
                      / %x9 ; ping
                      / %xA ; pong
                      / %xB-F ; reserved for further control
                              ; frames
                              ; 4 bits in length

frame-masked      = %x0
                  ; frame is not masked, no frame-masking-key
                  / %x1
                  ; frame is masked, frame-masking-key present
                  ; 1 bit in length

frame-payload-length = ( %x00-7D )
                      / ( %x7E frame-payload-length-16 )
                      / ( %x7F frame-payload-length-63 )
                      ; 7, 7+16, or 7+64 bits in length,
                      ; respectively

```



### 5.3. Client-to-Server Masking

A masked frame **MUST** have the field `frame-masked` set to 1, as defined in [Section 5.2](#).

The masking key is contained completely within the frame, as defined in [Section 5.2](#) as `frame-masking-key`. It is used to mask the "Payload data" defined in the same section as `frame-payload-data`, which includes "Extension data" and "Application data".

The masking key is a 32-bit value chosen at random by the client. When preparing a masked frame, the client **MUST** pick a fresh masking key from the set of allowed 32-bit values. The masking key needs to be unpredictable; thus, the masking key **MUST** be derived from a strong source of entropy, and the masking key for a given frame **MUST NOT** make it simple for a server/proxy to predict the masking key for a subsequent frame. The unpredictability of the masking key is essential to prevent authors of malicious applications from selecting the bytes that appear on the wire. RFC 4086 [[RFC4086](#)] discusses what entails a suitable source of entropy for security-sensitive applications.

The masking does not affect the length of the "Payload data". To convert masked data into unmasked data, or vice versa, the following algorithm is applied. The same algorithm applies regardless of the direction of the translation, e.g., the same steps are applied to mask the data as to unmask the data.

Octet *i* of the transformed data ("`transformed-octet-i`") is the XOR of octet *i* of the original data ("`original-octet-i`") with octet at index *i* modulo 4 of the masking key ("`masking-key-octet-j`"):

$$j = i \text{ MOD } 4$$

$$\text{transformed-octet-}i = \text{original-octet-}i \text{ XOR } \text{masking-key-octet-}j$$

The payload length, indicated in the framing as `frame-payload-length`, does **NOT** include the length of the masking key. It is the length of the "Payload data", e.g., the number of bytes following the masking key.

### 5.4. Fragmentation

The primary purpose of fragmentation is to allow sending a message that is of unknown size when the message is started without having to buffer that message. If messages couldn't be fragmented, then an endpoint would have to buffer the entire message so its length could be counted before the first byte is sent. With fragmentation, a server or intermediary may choose a reasonable size buffer and, when the buffer is full, write a fragment to the network.

A secondary use-case for fragmentation is for multiplexing, where it is not desirable for a large message on one logical channel to monopolize the output channel, so the multiplexing needs to be free to split the message into smaller fragments to better share the output channel. (Note that the multiplexing extension is not described in this document.)

Unless specified otherwise by an extension, frames have no semantic meaning. An intermediary might coalesce and/or split frames, if no extensions were negotiated by the client and the server or if some extensions were negotiated, but the intermediary understood all the extensions negotiated and knows how to coalesce and/or split frames in the presence of these extensions. One implication of this is that in absence of extensions, senders and receivers must not depend on the presence of specific frame boundaries.

The following rules apply to fragmentation:

- An unfragmented message consists of a single frame with the FIN bit set ([Section 5.2](#)) and an opcode other than 0.
- A fragmented message consists of a single frame with the FIN bit clear and an opcode other than 0, followed by zero or more frames with the FIN bit clear and the opcode set to 0, and terminated by a single frame with the FIN bit set and an opcode of 0. A fragmented message is conceptually equivalent to a single

larger message whose payload is equal to the concatenation of the payloads of the fragments in order; however, in the presence of extensions, this may not hold true as the extension defines the interpretation of the "Extension data" present. For instance, "Extension data" may only be present at the beginning of the first fragment and apply to subsequent fragments, or there may be "Extension data" present in each of the fragments that applies only to that particular fragment. In the absence of "Extension data", the following example demonstrates how fragmentation works.

EXAMPLE: For a text message sent as three fragments, the first fragment would have an opcode of 0x1 and a FIN bit clear, the second fragment would have an opcode of 0x0 and a FIN bit clear, and the third fragment would have an opcode of 0x0 and a FIN bit that is set.

- Control frames (see [Section 5.5](#)) MAY be injected in the middle of a fragmented message. Control frames themselves MUST NOT be fragmented.
- Message fragments MUST be delivered to the recipient in the order sent by the sender.
- The fragments of one message MUST NOT be interleaved between the fragments of another message unless an extension has been negotiated that can interpret the interleaving.
- An endpoint MUST be capable of handling control frames in the middle of a fragmented message.
- A sender MAY create fragments of any size for non-control messages.
- Clients and servers MUST support receiving both fragmented and unfragmented messages.
- As control frames cannot be fragmented, an intermediary MUST NOT attempt to change the fragmentation of a control frame.
- An intermediary MUST NOT change the fragmentation of a message if any reserved bit values are used and the meaning of these values is not known to the intermediary.
- An intermediary MUST NOT change the fragmentation of any message in the context of a connection where extensions have been negotiated and the intermediary is not aware of the semantics of the negotiated extensions. Similarly, an intermediary that didn't see the WebSocket handshake (and wasn't notified about its content) that resulted in a WebSocket connection MUST NOT change the fragmentation of any message of such connection.
- As a consequence of these rules, all fragments of a message are of the same type, as set by the first fragment's opcode. Since control frames cannot be fragmented, the type for all fragments in a message MUST be either text, binary, or one of the reserved opcodes.

NOTE: If control frames could not be interjected, the latency of a ping, for example, would be very long if behind a large message. Hence, the requirement of handling control frames in the middle of a fragmented message.

IMPLEMENTATION NOTE: In the absence of any extension, a receiver doesn't have to buffer the whole frame in order to process it. For example, if a streaming API is used, a part of a frame can be delivered to the application. However, note that this assumption might not hold true for all future WebSocket extensions.

## 5.5. Control Frames

Control frames are identified by opcodes where the most significant bit of the opcode is 1. Currently defined opcodes for control frames include 0x8 (Close), 0x9 (Ping), and 0xA (Pong). Opcodes 0xB-0xF are reserved for further control frames yet to be defined.

Control frames are used to communicate state about the WebSocket. Control frames can be interjected in the middle of a fragmented message.

All control frames MUST have a payload length of 125 bytes or less and MUST NOT be fragmented.

### 5.5.1. Close

The Close frame contains an opcode of 0x8.

The Close frame MAY contain a body (the "Application data" portion of the frame) that indicates a reason for closing, such as an endpoint shutting down, an endpoint having received a frame too large, or an endpoint

having received a frame that does not conform to the format expected by the endpoint. If there is a body, the first two bytes of the body **MUST** be a 2-byte unsigned integer (in network byte order) representing a status code with value `/code/` defined in [Section 7.4](#). Following the 2-byte integer, the body **MAY** contain UTF#8-encoded data with value `/reason/`, the interpretation of which is not defined by this specification. This data is not necessarily human readable but may be useful for debugging or passing information relevant to the script that opened the connection. As the data is not guaranteed to be human readable, clients **MUST NOT** show it to end users.

Close frames sent from client to server must be masked as per [Section 5.3](#).

The application **MUST NOT** send any more data frames after sending a Close frame.

If an endpoint receives a Close frame and did not previously send a Close frame, the endpoint **MUST** send a Close frame in response. (When sending a Close frame in response, the endpoint typically echos the status code it received.) It **SHOULD** do so as soon as practical. An endpoint **MAY** delay sending a Close frame until its current message is sent (for instance, if the majority of a fragmented message is already sent, an endpoint **MAY** send the remaining fragments before sending a Close frame). However, there is no guarantee that the endpoint that has already sent a Close frame will continue to process data.

After both sending and receiving a Close message, an endpoint considers the WebSocket connection closed and **MUST** close the underlying TCP connection. The server **MUST** close the underlying TCP connection immediately; the client **SHOULD** wait for the server to close the connection but **MAY** close the connection at any time after sending and receiving a Close message, e.g., if it has not received a TCP Close from the server in a reasonable time period.

If a client and server both send a Close message at the same time, both endpoints will have sent and received a Close message and should consider the WebSocket connection closed and close the underlying TCP connection.

### 5.5.2. Ping

The Ping frame contains an opcode of 0x9.

A Ping frame **MAY** include "Application data".

Upon receipt of a Ping frame, an endpoint **MUST** send a Pong frame in response, unless it already received a Close frame. It **SHOULD** respond with Pong frame as soon as is practical. Pong frames are discussed in [Section 5.5.3](#).

An endpoint **MAY** send a Ping frame any time after the connection is established and before the connection is closed.

NOTE: A Ping frame may serve either as a keepalive or as a means to verify that the remote endpoint is still responsive.

### 5.5.3. Pong

The Pong frame contains an opcode of 0xA.

[Section 5.5.2](#) details requirements that apply to both Ping and Pong frames.

A Pong frame sent in response to a Ping frame must have identical "Application data" as found in the message body of the Ping frame being replied to.

If an endpoint receives a Ping frame and has not yet sent Pong frame(s) in response to previous Ping frame(s), the endpoint **MAY** elect to send a Pong frame for only the most recently processed Ping frame.

A Pong frame **MAY** be sent unsolicited. This serves as a unidirectional heartbeat. A response to an unsolicited Pong frame is not expected.

## 5.6. Data Frames

Data frames (e.g., non-control frames) are identified by opcodes where the most significant bit of the opcode is 0. Currently defined opcodes for data frames include 0x1 (Text), 0x2 (Binary). Opcodes 0x3-0x7 are reserved for further non-control frames yet to be defined.

Data frames carry application-layer and/or extension-layer data. The opcode determines the interpretation of the data:

- |        |  |
|--------|--|
| Text   | The "Payload data" is text data encoded as UTF-8. Note that a particular text frame might include a partial UTF-8 sequence; however, the whole message <b>MUST</b> contain valid UTF-8. Invalid UTF-8 in reassembled messages is handled as described in <a href="#">Section 8.1</a> . |
| Binary | The "Payload data" is arbitrary binary data whose interpretation is solely up to the application layer.  |

## 5.7. Examples

- A single-frame unmasked text message
  - 0x81 0x05 0x48 0x65 0x6c 0x6c 0x6f (contains "Hello")
- A single-frame masked text message
  - 0x81 0x85 0x37 0xfa 0x21 0x3d 0x7f 0x9f 0x4d 0x51 0x58 (contains "Hello")
- A fragmented unmasked text message
  - 0x01 0x03 0x48 0x65 0x6c (contains "Hel")
  - 0x80 0x02 0x6c 0x6f (contains "lo")
- Unmasked Ping request and masked Ping response
  - 0x89 0x05 0x48 0x65 0x6c 0x6c 0x6f (contains a body of "Hello", but the contents of the body are arbitrary)
  - 0x8a 0x85 0x37 0xfa 0x21 0x3d 0x7f 0x9f 0x4d 0x51 0x58 (contains a body of "Hello", matching the body of the ping)
- 256 bytes binary message in a single unmasked frame
  - 0x82 0x7E 0x0100 [256 bytes of binary data]
- 64KiB binary message in a single unmasked frame
  - 0x82 0x7F 0x0000000000010000 [65536 bytes of binary data]

## 5.8. Extensibility

The protocol is designed to allow for extensions, which will add capabilities to the base protocol. The endpoints of a connection **MUST** negotiate the use of any extensions during the opening handshake. This specification provides opcodes 0x3 through 0x7 and 0xB through 0xF, the "Extension data" field, and the frame-rsv1, frame-rsv2, and frame-rsv3 bits of the frame header for use by extensions. The negotiation of extensions is discussed in further detail in [Section 9.1](#). Below are some anticipated uses of extensions. This list is neither complete nor prescriptive.

- "Extension data" may be placed in the "Payload data" before the "Application data".
- Reserved bits can be allocated for per-frame needs.
- Reserved opcode values can be defined.
- Reserved bits can be allocated to the opcode field if more opcode values are needed.
- A reserved bit or an "extension" opcode can be defined that allocates additional bits out of the "Payload data" to define larger opcodes or more per-frame bits.

## 6. Sending and Receiving Data

### 6.1. Sending Data

To *Send a WebSocket Message* comprising of /data/ over a WebSocket connection, an endpoint **MUST** perform the following steps.

1. The endpoint **MUST** ensure the WebSocket connection is in the OPEN state (cf. Sections 4.1 and 4.2.2.) If at any point the state of the WebSocket connection changes, the endpoint **MUST** abort the following steps.
2. An endpoint **MUST** encapsulate the /data/ in a WebSocket frame as defined in Section 5.2. If the data to be sent is large or if the data is not available in its entirety at the point the endpoint wishes to begin sending the data, the endpoint **MAY** alternately encapsulate the data in a series of frames as defined in Section 5.4.
3. The opcode (frame-opcode) of the first frame containing the data **MUST** be set to the appropriate value from Section 5.2 for data that is to be interpreted by the recipient as text or binary data.
4. The FIN bit (frame-fin) of the last frame containing the data **MUST** be set to 1 as defined in Section 5.2.
5. If the data is being sent by the client, the frame(s) **MUST** be masked as defined in Section 5.3.
6. If any extensions (Section 9) have been negotiated for the WebSocket connection, additional considerations may apply as per the definition of those extensions.
7. The frame(s) that have been formed **MUST** be transmitted over the underlying network connection.

### 6.2. Receiving Data

To receive WebSocket data, an endpoint listens on the underlying network connection. Incoming data **MUST** be parsed as WebSocket frames as defined in Section 5.2. If a control frame (Section 5.5) is received, the frame **MUST** be handled as defined by Section 5.5. Upon receiving a data frame (Section 5.6), the endpoint **MUST** note the /type/ of the data as defined by the opcode (frame-opcode) from Section 5.2. The "Application data" from this frame is defined as the /data/ of the message. If the frame comprises an unfragmented message (Section 5.4), it is said that *A WebSocket Message Has Been Received* with type /type/ and data /data/. If the frame is part of a fragmented message, the "Application data" of the subsequent data frames is concatenated to form the /data/. When the last fragment is received as indicated by the FIN bit (frame-fin), it is said that *A WebSocket Message Has Been Received* with data /data/ (comprised of the concatenation of the "Application data" of the fragments) and type /type/ (noted from the first frame of the fragmented message). Subsequent data frames **MUST** be interpreted as belonging to a new WebSocket message.

Extensions (Section 9) **MAY** change the semantics of how data is read, specifically including what comprises a message boundary. Extensions, in addition to adding "Extension data" before the "Application data" in a payload, **MAY** also modify the "Application data" (such as by compressing it).

A server **MUST** remove masking for data frames received from a client as described in Section 5.3.

## 7. Closing the Connection

### 7.1. Definitions

#### 7.1.1. Close the WebSocket Connection

To *Close the WebSocket Connection*, an endpoint closes the underlying TCP connection. An endpoint SHOULD use a method that cleanly closes the TCP connection, as well as the TLS session, if applicable, discarding any trailing bytes that may have been received. An endpoint MAY close the connection via any means available when necessary, such as when under attack.

The underlying TCP connection, in most normal cases, SHOULD be closed first by the server, so that it holds the TIME\_WAIT state and not the client (as this would prevent it from re-opening the connection for 2 maximum segment lifetimes (2MSL), while there is no corresponding server impact as a TIME\_WAIT connection is immediately reopened upon a new SYN with a higher seq number). In abnormal cases (such as not having received a TCP Close from the server after a reasonable amount of time) a client MAY initiate the TCP Close. As such, when a server is instructed to *Close the WebSocket Connection* it SHOULD initiate a TCP Close immediately, and when a client is instructed to do the same, it SHOULD wait for a TCP Close from the server.

As an example of how to obtain a clean closure in C using Berkeley sockets, one would call shutdown() with SHUT\_WR on the socket, call recv() until obtaining a return value of 0 indicating that the peer has also performed an orderly shutdown, and finally call close() on the socket.

#### 7.1.2. Start the WebSocket Closing Handshake

To *Start the WebSocket Closing Handshake* with a status code (Section 7.4) /code/ and an optional close reason (Section 7.1.6) /reason/, an endpoint MUST send a Close control frame, as described in Section 5.5.1, whose status code is set to /code/ and whose close reason is set to /reason/. Once an endpoint has both sent and received a Close control frame, that endpoint SHOULD *Close the WebSocket Connection* as defined in Section 7.1.1.

#### 7.1.3. The WebSocket Closing Handshake is Started

Upon either sending or receiving a Close control frame, it is said that *The WebSocket Closing Handshake is Started* and that the WebSocket connection is in the CLOSING state.

#### 7.1.4. The WebSocket Connection is Closed

When the underlying TCP connection is closed, it is said that *The WebSocket Connection is Closed* and that the WebSocket connection is in the CLOSED state. If the TCP connection was closed after the WebSocket closing handshake was completed, the WebSocket connection is said to have been closed *cleanly*.

If the WebSocket connection could not be established, it is also said that *The WebSocket Connection is Closed*, but not *cleanly*.

#### 7.1.5. The WebSocket Connection Close Code

As defined in Sections 5.5.1 and 7.4, a Close control frame may contain a status code indicating a reason for closure. A closing of the WebSocket connection may be initiated by either endpoint, potentially simultaneously. *The WebSocket Connection Close Code* is defined as the status code (Section 7.4) contained in the first Close control frame received by the application implementing this protocol. If this Close control frame contains no status code, *The WebSocket Connection Close Code* is considered to be 1005. If *The WebSocket Connection is Closed* and no Close control frame was received by the endpoint (such as could occur if the underlying transport connection is lost), *The WebSocket Connection Close Code* is considered to be 1006.

NOTE: Two endpoints may not agree on the value of *The WebSocket Connection Close Code*. As an example, if the remote endpoint sent a Close frame but the local application has not yet read the data containing the

Close frame from its socket's receive buffer, and the local application independently decided to close the connection and send a Close frame, both endpoints will have sent and received a Close frame and will not send further Close frames. Each endpoint will see the status code sent by the other end as *The WebSocket Connection Close Code*. As such, it is possible that the two endpoints may not agree on the value of *The WebSocket Connection Close Code* in the case that both endpoints *Start the WebSocket Closing Handshake* independently and at roughly the same time.

### 7.1.6. The WebSocket Connection Close Reason

As defined in Sections 5.5.1 and 7.4, a Close control frame may contain a status code indicating a reason for closure, followed by UTF#8-encoded data, the interpretation of said data being left to the endpoints and not defined by this protocol. A closing of the WebSocket connection may be initiated by either endpoint, potentially simultaneously. *The WebSocket Connection Close Reason* is defined as the UTF#8-encoded data following the status code (Section 7.4) contained in the first Close control frame received by the application implementing this protocol. If there is no such data in the Close control frame, *The WebSocket Connection Close Reason* is the empty string.

NOTE: Following the same logic as noted in Section 7.1.5, two endpoints may not agree on *The WebSocket Connection Close Reason*.

### 7.1.7. Fail the WebSocket Connection

Certain algorithms and specifications require an endpoint to *Fail the WebSocket Connection*. To do so, the client **MUST** *Close the WebSocket Connection*, and **MAY** report the problem to the user (which would be especially useful for developers) in an appropriate manner. Similarly, to do so, the server **MUST** *Close the WebSocket Connection*, and **SHOULD** log the problem.

If *The WebSocket Connection is Established* prior to the point where the endpoint is required to *Fail the WebSocket Connection*, the endpoint **SHOULD** send a Close frame with an appropriate status code (Section 7.4) before proceeding to *Close the WebSocket Connection*. An endpoint **MAY** omit sending a Close frame if it believes the other side is unlikely to be able to receive and process the Close frame, due to the nature of the error that led the WebSocket connection to fail in the first place. An endpoint **MUST NOT** continue to attempt to process data (including a responding Close frame) from the remote endpoint after being instructed to *Fail the WebSocket Connection*.

Except as indicated above or as specified by the application layer (e.g., a script using the WebSocket API), clients **SHOULD NOT** close the connection.

## 7.2. Abnormal Closures

### 7.2.1. Client-Initiated Closure

Certain algorithms, in particular during the opening handshake, require the client to *Fail the WebSocket Connection*. To do so, the client **MUST** *Fail the WebSocket Connection* as defined in Section 7.1.7.

If at any point the underlying transport layer connection is unexpectedly lost, the client **MUST** *Fail the WebSocket Connection*.

Except as indicated above or as specified by the application layer (e.g., a script using the WebSocket API), clients **SHOULD NOT** close the connection.

### 7.2.2. Server-Initiated Closure

Certain algorithms require or recommend that the server *Abort the WebSocket Connection* during the opening handshake. To do so, the server **MUST** simply *Close the WebSocket Connection* (Section 7.1.1).

### 7.2.3. Recovering from Abnormal Closure

Abnormal closures may be caused by any number of reasons. Such closures could be the result of a transient error, in which case reconnecting may lead to a good connection and a resumption of normal operations. Such closures may also be the result of a nontransient problem, in which case if each deployed client experiences an abnormal closure and immediately and persistently tries to reconnect, the server may experience what amounts to a denial-of-service attack by a large number of clients trying to reconnect. The end result of such a scenario could be that the service is unable to recover in a timely manner or recovery is made much more difficult.

To prevent this, clients **SHOULD** use some form of backoff when trying to reconnect after abnormal closures as described in this section.

The first reconnect attempt **SHOULD** be delayed by a random amount of time. The parameters by which this random delay is chosen are left to the client to decide; a value chosen randomly between 0 and 5 seconds is a reasonable initial delay though clients **MAY** choose a different interval from which to select a delay length based on implementation experience and particular application.

Should the first reconnect attempt fail, subsequent reconnect attempts **SHOULD** be delayed by increasingly longer amounts of time, using a method such as truncated binary exponential backoff.

## 7.3. Normal Closure of Connections

Servers **MAY** close the WebSocket connection whenever desired. Clients **SHOULD NOT** close the WebSocket connection arbitrarily. In either case, an endpoint initiates a closure by following the procedures to *Start the WebSocket Closing Handshake* (Section 7.1.2).

## 7.4. Status Codes

When closing an established connection (e.g., when sending a Close frame, after the opening handshake has completed), an endpoint **MAY** indicate a reason for closure. The interpretation of this reason by an endpoint, and the action an endpoint should take given this reason, are left undefined by this specification. This specification defines a set of pre-defined status codes and specifies which ranges may be used by extensions, frameworks, and end applications. The status code and any associated textual message are optional components of a Close frame.

### 7.4.1. Defined Status Codes

Endpoints **MAY** use the following pre-defined status codes when sending a Close frame.

- 1000 1000 indicates a normal closure, meaning that the purpose for which the connection was established has been fulfilled.
- 1001 1001 indicates that an endpoint is "going away", such as a server going down or a browser having navigated away from a page.
- 1002 1002 indicates that an endpoint is terminating the connection due to a protocol error.
- 1003 1003 indicates that an endpoint is terminating the connection because it has received a type of data it cannot accept (e.g., an endpoint that understands only text data **MAY** send this if it receives a binary message).
- 1004 Reserved. The specific meaning might be defined in the future.
- 1005 1005 is a reserved value and **MUST NOT** be set as a status code in a Close control frame by an endpoint. It is designated for use in applications expecting a status code to indicate that no status code was actually present.
- 1006 1006 is a reserved value and **MUST NOT** be set as a status code in a Close control frame by an endpoint. It is designated for use in applications expecting a status code to indicate that the connection was closed abnormally, e.g., without sending or receiving a Close control frame.
- 1007



- 1007 indicates that an endpoint is terminating the connection because it has received data within a message that was not consistent with the type of the message (e.g., non-UTF-8 [RFC3629] data within a text message).
- 1008 1008 indicates that an endpoint is terminating the connection because it has received a message that violates its policy. This is a generic status code that can be returned when there is no other more suitable status code (e.g., 1003 or 1009) or if there is a need to hide specific details about the policy.
- 1009 1009 indicates that an endpoint is terminating the connection because it has received a message that is too big for it to process.
- 1010 1010 indicates that an endpoint (client) is terminating the connection because it has expected the server to negotiate one or more extension, but the server didn't return them in the response message of the WebSocket handshake. The list of extensions that are needed SHOULD appear in the /reason/ part of the Close frame. Note that this status code is not used by the server, because it can fail the WebSocket handshake instead.
- 1011 1011 indicates that a server is terminating the connection because it encountered an unexpected condition that prevented it from fulfilling the request.
- 1015 1015 is a reserved value and MUST NOT be set as a status code in a Close control frame by an endpoint. It is designated for use in applications expecting a status code to indicate that the connection was closed due to a failure to perform a TLS handshake (e.g., the server certificate can't be verified).

#### 7.4.2. Reserved Status Code Ranges

- 0-999 Status codes in the range 0-999 are not used.
- 1000-2999 Status codes in the range 1000-2999 are reserved for definition by this protocol, its future revisions, and extensions specified in a permanent and readily available public specification.
- 3000-3999 Status codes in the range 3000-3999 are reserved for use by libraries, frameworks, and applications. These status codes are registered directly with IANA. The interpretation of these codes is undefined by this protocol.
- 4000-4999 Status codes in the range 4000-4999 are reserved for private use and thus can't be registered. Such codes can be used by prior agreements between WebSocket applications. The interpretation of these codes is undefined by this protocol.

## 8. Error Handling

### 8.1. Handling Errors in UTF-8-Encoded Data

When an endpoint is to interpret a byte stream as UTF-8 but finds that the byte stream is not, in fact, a valid UTF-8 stream, that endpoint **MUST** *Fail the WebSocket Connection*. This rule applies both during the opening handshake and during subsequent data exchange.

## 9. Extensions

WebSocket clients MAY request extensions to this specification, and WebSocket servers MAY accept some or all extensions requested by the client. A server MUST NOT respond with any extension not requested by the client. If extension parameters are included in negotiations between the client and the server, those parameters MUST be chosen in accordance with the specification of the extension to which the parameters apply.

### 9.1. Negotiating Extensions

A client requests extensions by including a `|Sec-WebSocket-Extensions|` header field, which follows the normal rules for HTTP header fields (see [RFC2616], Section 4.2) and the value of the header field is defined by the following ABNF [RFC2616]. Note that this section is using ABNF syntax/rules from [RFC2616], including the "implied \*LWS rule". If a value is received by either the client or the server during negotiation that does not conform to the ABNF below, the recipient of such malformed data MUST immediately *Fail the WebSocket Connection*.

```
Sec-WebSocket-Extensions = extension-list
extension-list = 1#extension
extension = extension-token *( ";" extension-param )
extension-token = registered-token
registered-token = token
extension-param = token [ "=" (token | quoted-string) ]
    ;When using the quoted-string syntax variant, the value
    ;after quoted-string unescaping MUST conform to the
    ;'token' ABNF.
```

Note that like other HTTP header fields, this header field MAY be split or combined across multiple lines. Ergo, the following are equivalent:

```
Sec-WebSocket-Extensions: foo
Sec-WebSocket-Extensions: bar; baz=2
```

is exactly equivalent to

```
Sec-WebSocket-Extensions: foo, bar; baz=2
```

Any extension-token used MUST be a registered token (see Section 11.4). The parameters supplied with any given extension MUST be defined for that extension. Note that the client is only offering to use any advertised extensions and MUST NOT use them unless the server indicates that it wishes to use the extension.

Note that the order of extensions is significant. Any interactions between multiple extensions MAY be defined in the documents defining the extensions. In the absence of such definitions, the interpretation is that the header fields listed by the client in its request represent a preference of the header fields it wishes to use, with the first options listed being most preferable. The extensions listed by the server in response represent the extensions actually in use for the connection. Should the extensions modify the data and/or framing, the order of operations on the data should be assumed to be the same as the order in which the extensions are listed in the server's response in the opening handshake.

For example, if there are two extensions "foo" and "bar" and if the header field `|Sec-WebSocket-Extensions|` sent by the server has the value "foo, bar", then operations on the data will be made as `bar(foo(data))`, be those changes to the data itself (such as compression) or changes to the framing that may "stack".

Non-normative examples of acceptable extension header fields (note that long lines are folded for readability):

```
Sec-WebSocket-Extensions: deflate-stream
Sec-WebSocket-Extensions: mux; max-channels=4; flow-control,
  deflate-stream
Sec-WebSocket-Extensions: private-extension
```

A server accepts one or more extensions by including a |Sec-WebSocket-Extensions| header field containing one or more extensions that were requested by the client. The interpretation of any extension parameters, and what constitutes a valid response by a server to a requested set of parameters by a client, will be defined by each such extension.

## 9.2. Known Extensions

Extensions provide a mechanism for implementations to opt-in to additional protocol features. This document doesn't define any extension, but implementations MAY use extensions defined separately.

## 10. Security Considerations

This section describes some security considerations applicable to the WebSocket Protocol. Specific security considerations are described in subsections of this section.

### 10.1. Non-Browser Clients

The WebSocket Protocol protects against malicious JavaScript running inside a trusted application such as a web browser, for example, by checking of the |Origin| header field (see below). See [Section 1.6](#) for additional details. Such assumptions don't hold true in the case of a more-capable client.

While this protocol is intended to be used by scripts in web pages, it can also be used directly by hosts. Such hosts are acting on their own behalf and can therefore send fake |Origin| header fields, misleading the server. Servers should therefore be careful about assuming that they are talking directly to scripts from known origins and must consider that they might be accessed in unexpected ways. In particular, a server should not trust that any input is valid.

EXAMPLE: If the server uses input as part of SQL queries, all input text should be escaped before being passed to the SQL server, lest the server be susceptible to SQL injection.

### 10.2. Origin Considerations

Servers that are not intended to process input from any web page but only for certain sites SHOULD verify the |Origin| field is an origin they expect. If the origin indicated is unacceptable to the server, then it SHOULD respond to the WebSocket handshake with a reply containing HTTP 403 Forbidden status code.

The |Origin| header field protects from the attack cases when the untrusted party is typically the author of a JavaScript application that is executing in the context of the trusted client. The client itself can contact the server and, via the mechanism of the |Origin| header field, determine whether to extend those communication privileges to the JavaScript application. The intent is not to prevent non-browsers from establishing connections but rather to ensure that trusted browsers under the control of potentially malicious JavaScript cannot fake a WebSocket handshake.

### 10.3. Attacks On Infrastructure (Masking)

In addition to endpoints being the target of attacks via WebSockets, other parts of web infrastructure, such as proxies, may be the subject of an attack.

As this protocol was being developed, an experiment was conducted to demonstrate a class of attacks on proxies that led to the poisoning of caching proxies deployed in the wild [TALKING]. The general form of the attack was to establish a connection to a server under the "attacker's" control, perform an UPGRADE on the HTTP connection similar to what the WebSocket Protocol does to establish a connection, and subsequently send data over that UPGRADED connection that looked like a GET request for a specific known resource (which in an attack would likely be something like a widely deployed script for tracking hits or a resource on an ad-serving network). The remote server would respond with something that looked like a response to the fake GET request, and this response would be cached by a nonzero percentage of deployed intermediaries, thus poisoning the cache. The net effect of this attack would be that if a user could be convinced to visit a website the attacker controlled, the attacker could potentially poison the cache for that user and other users behind the same cache and run malicious script on other origins, compromising the web security model.

To avoid such attacks on deployed intermediaries, it is not sufficient to prefix application-supplied data with framing that is not compliant with HTTP, as it is not possible to exhaustively discover and test that each nonconformant intermediary does not skip such non-HTTP framing and act incorrectly on the frame payload. Thus, the defense adopted is to mask all data from the client to the server, so that the remote script (attacker) does not have control over how the data being sent appears on the wire and thus cannot construct a message that could be misinterpreted by an intermediary as an HTTP request.

Clients **MUST** choose a new masking key for each frame, using an algorithm that cannot be predicted by end applications that provide data. For example, each masking could be drawn from a cryptographically strong random number generator. If the same key is used or a decipherable pattern exists for how the next key is chosen, the attacker can send a message that, when masked, could appear to be an HTTP request (by taking the message the attacker wishes to see on the wire and masking it with the next masking key to be used, the masking key will effectively unmask the data when the client applies it).

It is also necessary that once the transmission of a frame from a client has begun, the payload (application-supplied data) of that frame must not be capable of being modified by the application. Otherwise, an attacker could send a long frame where the initial data was a known value (such as all zeros), compute the masking key being used upon receipt of the first part of the data, and then modify the data that is yet to be sent in the frame to appear as an HTTP request when masked. (This is essentially the same problem described in the previous paragraph with using a known or predictable masking key.) If additional data is to be sent or data to be sent is somehow changed, that new or changed data must be sent in a new frame and thus with a new masking key. In short, once transmission of a frame begins, the contents must not be modifiable by the remote script (application).

The threat model being protected against is one in which the client sends data that appears to be an HTTP request. As such, the channel that needs to be masked is the data from the client to the server. The data from the server to the client can be made to look like a response, but to accomplish this request, the client must also be able to forge a request. As such, it was not deemed necessary to mask data in both directions (the data from the server to the client is not masked).

Despite the protection provided by masking, non-compliant HTTP proxies will still be vulnerable to poisoning attacks of this type by clients and servers that do not apply masking.

#### 10.4. Implementation-Specific Limits

Implementations that have implementation- and/or platform-specific limitations regarding the frame size or total message size after reassembly from multiple frames **MUST** protect themselves against exceeding those limits. (For example, a malicious endpoint can try to exhaust its peer's memory or mount a denial-of-service attack by sending either a single big frame (e.g., of size  $2^{60}$ ) or by sending a long stream of small frames that are a part of a fragmented message.) Such an implementation **SHOULD** impose a limit on frame sizes and the total message size after reassembly from multiple frames.

#### 10.5. WebSocket Client Authentication

This protocol doesn't prescribe any particular way that servers can authenticate clients during the WebSocket handshake. The WebSocket server can use any client authentication mechanism available to a generic HTTP server, such as cookies, HTTP authentication, or TLS authentication.

#### 10.6. Connection Confidentiality and Integrity

Connection confidentiality and integrity is provided by running the WebSocket Protocol over TLS (wss URIs). WebSocket implementations **MUST** support TLS and **SHOULD** employ it when communicating with their peers.

For connections using TLS, the amount of benefit provided by TLS depends greatly on the strength of the algorithms negotiated during the TLS handshake. For example, some TLS cipher mechanisms don't provide connection confidentiality. To achieve reasonable levels of protection, clients should use only Strong TLS algorithms. "Web Security Context: User Interface Guidelines" [W3C.REC-wsc-ui-20100812] discusses what constitutes Strong TLS algorithms. [RFC5246] provides additional guidance in Appendix A.5 and Appendix D.3.

## 10.7. Handling of Invalid Data

Incoming data **MUST** always be validated by both clients and servers. If, at any time, an endpoint is faced with data that it does not understand or that violates some criteria by which the endpoint determines safety of input, or when the endpoint sees an opening handshake that does not correspond to the values it is expecting (e.g., incorrect path or origin in the client request), the endpoint **MAY** drop the TCP connection. If the invalid data was received after a successful WebSocket handshake, the endpoint **SHOULD** send a Close frame with an appropriate status code ([Section 7.4](#)) before proceeding to *Close the WebSocket Connection*. Use of a Close frame with an appropriate status code can help in diagnosing the problem. If the invalid data is sent during the WebSocket handshake, the server **SHOULD** return an appropriate HTTP [\[RFC2616\]](#) status code.

A common class of security problems arises when sending text data using the wrong encoding. This protocol specifies that messages with a Text data type (as opposed to Binary or other types) contain UTF#8-encoded data. Although the length is still indicated and applications implementing this protocol should use the length to determine where the frame actually ends, sending data in an improper encoding may still break assumptions that applications built on top of this protocol may make, leading to anything from misinterpretation of data to loss of data or potential security bugs.

## 10.8. Use of SHA-1 by the WebSocket Handshake

The WebSocket handshake described in this document doesn't depend on any security properties of SHA-1, such as collision resistance or resistance to the second pre-image attack (as described in [\[RFC4270\]](#)).

## 11. IANA Considerations

### 11.1. Registration of New URI Schemes

#### 11.1.1. Registration of "ws" Scheme

A |ws| URI identifies a WebSocket server and resource name.

URI scheme name

ws

Status

Permanent

URI scheme syntax

Using the ABNF [\[RFC5234\]](#) syntax and ABNF terminals from the URI specification [\[RFC3986\]](#):

```
"ws:"
"//"
authority
path-
abempty
[ "?"
query ]
```

The <path-abempty> and <query> [\[RFC3986\]](#) components form the resource name sent to the server to identify the kind of service desired. Other components



URI scheme semantics

have the meanings described in [\[RFC3986\]](#).

The only operation for this scheme is to open a connection using the WebSocket Protocol.

Encoding considerations

Characters in the host component that are excluded by the syntax defined above MUST be converted from Unicode to ASCII as specified in [\[RFC3987\]](#) or its replacement. For the purposes of scheme-based normalization, Internationalized Domain Name (IDN) forms of the

host component and their conversions to punycode are considered equivalent (see Section 5.3.3 of [\[RFC3987\]](#)). Characters in other components that are excluded by the syntax defined above MUST be converted from Unicode to ASCII by first encoding the characters as UTF-8 and then replacing the corresponding bytes using their percent-encoded form as defined in the URI [\[RFC3986\]](#) and Internationalized Resource

	Identifier (IRI) <a href="#">[RFC3987]</a> specifications.
Applications/protocols that use this URI scheme name	WebSocket Protocol
Interoperability considerations	Use of WebSocket requires use of HTTP version 1.1 or higher.
Security considerations	See "Security Considerations" section.
Contact	HYBI WG <hybi@ietf.org>
Author/Change controller	IETF <iesg@ietf.org>
References	RFC 6455

### 11.1.2. Registration of "wss" Scheme

A [wss] URI identifies a WebSocket server and resource name and indicates that traffic over that connection is to be protected via TLS (including standard benefits of TLS such as data confidentiality and integrity and endpoint authentication).

URI scheme name	wss
Status	Permanent
URI scheme syntax	Using the ABNF <a href="#">[RFC5234]</a> syntax and ABNF terminals from the URI specification <a href="#">[RFC3986]</a> :

```

"wss:"
"//"
authority
path-
abempty
[ "?"
query ]

```

The  
 <path-  
 abempty>  
 and  
 <query>  
 components  
 form  
 the  
 resource  
 name  
 sent  
 to the  
 server  
 to  
 identify  
 the  
 kind of  
 service  
 desired.  
 Other  
 components  
 have  
 the  
 meanings  
 described  
 in  
[\[RFC3986\]](#).

URI scheme semantics

The  
 only  
 operation  
 for this  
 scheme  
 is to  
 open a  
 connection  
 using  
 the  
 WebSocket  
 Protocol,  
 encrypted  
 using  
 TLS.

Encoding considerations

Characters in the host component that are excluded by the syntax defined above MUST be converted from Unicode to ASCII as specified in [\[RFC3987\]](#) or its replacement. For the purposes of scheme-based normalization IDN forms of the host component and their conversions to punycode are considered equivalent (see Section 5.3.3 of [\[RFC3987\]](#)). Characters in other components that are excluded by the syntax defined

	above MUST be converted from Unicode to ASCII by first encoding the characters as UTF-8 and then replacing the corresponding bytes using their percent- encoded form as defined in the URI <a href="#">[RFC3986]</a> and IRI <a href="#">[RFC3987]</a> specifications.
Applications/protocols that use this URI scheme name	WebSocket Protocol over TLS
Interoperability considerations	Use of WebSocket requires use of HTTP version 1.1 or higher.
Security considerations	See "Security Considerations" section.
Contact	HYBI WG <hybi@ietf.org>

Author/Change controller	IETF <iesg@ietf.org>
References	RFC 6455

## 11.2. Registration of the "WebSocket" HTTP Upgrade Keyword

This section defines a keyword registered in the HTTP Upgrade Tokens Registry as per RFC 2817 [RFC2817].

Name of token	WebSocket
Author/Change controller	IETF <iesg@ietf.org>
Contact	HYBI <hybi@ietf.org>
References	RFC 6455

## 11.3. Registration of New HTTP Header Fields

### 11.3.1. Sec-WebSocket-Key

This section describes a header field registered in the Permanent Message Header Field Names registry [RFC3864].

Header field name	Sec-WebSocket-Key
Applicable protocol	http
Status	standard
Author/Change controller	IETF
Specification document(s)	RFC 6455
Related information	This header field is only used for WebSocket opening handshake.

The |Sec-WebSocket-Key| header field is used in the WebSocket opening handshake. It is sent from the client to the server to provide part of the information used by the server to prove that it received a valid WebSocket opening handshake. This helps ensure that the server does not accept connections from non-WebSocket clients (e.g., HTTP clients) that are being abused to send data to unsuspecting WebSocket servers.

The |Sec-WebSocket-Key| header field **MUST NOT** appear more than once in an HTTP request.

### 11.3.2. Sec-WebSocket-Extensions

This section describes a header field for registration in the Permanent Message Header Field Names registry [RFC3864].

Header field name	Sec-WebSocket-Extensions
Applicable protocol	http
Status	standard
Author/Change controller	IETF
Specification document(s)	RFC 6455
Related information	This header field is only used for WebSocket opening handshake.

The |Sec-WebSocket-Extensions| header field is used in the WebSocket opening handshake. It is initially sent from the client to the server, and then subsequently sent from the server to the client, to agree on a set of protocol-level extensions to use for the duration of the connection.

The |Sec-WebSocket-Extensions| header field MAY appear multiple times in an HTTP request (which is logically the same as a single |Sec-WebSocket-Extensions| header field that contains all values). However, the |Sec-WebSocket-Extensions| header field MUST NOT appear more than once in an HTTP response.

### 11.3.3. Sec-WebSocket-Accept

This section describes a header field registered in the Permanent Message Header Field Names registry [\[RFC3864\]](#).

Header field name	Sec-WebSocket-Accept
Applicable protocol	http
Status	standard
Author/Change controller	IETF
Specification document(s)	RFC 6455
Related information	This header field is only used for the WebSocket opening handshake.

The |Sec-WebSocket-Accept| header field is used in the WebSocket opening handshake. It is sent from the server to the client to confirm that the server is willing to initiate the WebSocket connection.

The |Sec-WebSocket-Accept| header MUST NOT appear more than once in an HTTP response.

### 11.3.4. Sec-WebSocket-Protocol

This section describes a header field registered in the Permanent Message Header Field Names registry [\[RFC3864\]](#).

Header field name	Sec-WebSocket-Protocol
Applicable protocol	http
Status	standard
Author/Change controller	IETF
Specification document(s)	RFC 6455
Related information	This header field is only used for the WebSocket opening handshake.

The |Sec-WebSocket-Protocol| header field is used in the WebSocket opening handshake. It is sent from the client to the server and back from the server to the client to confirm the subprotocol of the connection. This enables scripts to both select a subprotocol and be sure that the server agreed to serve that subprotocol.

The |Sec-WebSocket-Protocol| header field MAY appear multiple times in an HTTP request (which is logically the same as a single |Sec-WebSocket-Protocol| header field that contains all values). However, the |Sec-WebSocket-Protocol| header field MUST NOT appear more than once in an HTTP response.

### 11.3.5. Sec-WebSocket-Version

This section describes a header field registered in the Permanent Message Header Field Names registry [\[RFC3864\]](#).

Header field name	Sec-WebSocket-Version
Applicable protocol	http
Status	standard
Author/Change controller	IETF
Specification document(s)	RFC 6455





A reference to the document in which the subprotocol being used with the WebSocket Protocol is defined.

WebSocket Subprotocol names are to be subject to the "First Come First Served" IANA registration policy [RFC5226].

## 11.6. WebSocket Version Number Registry

This specification creates a new IANA registry for WebSocket Version Numbers to be used with the WebSocket Protocol in accordance with the principles set out in RFC 5226 [RFC5226].

As part of this registry, IANA maintains the following information:

Version Number	The version number to be used in the [Sec-WebSocket-Version] is specified in <a href="#">Section 4.1</a> of this specification. The value must be a non-negative integer in the range between 0 and 255 (inclusive).
Reference	The RFC requesting a new version number or a draft name with version number (see below).
Status	Either "Interim" or "Standard". See below for description.

A version number is designated as either "Interim" or "Standard".

A "Standard" version number is documented in an RFC and used to identify a major, stable version of the WebSocket protocol, such as the version defined by this RFC. "Standard" version numbers are subject to the "IETF Review" IANA registration policy [RFC5226].

An "Interim" version number is documented in an Internet-Draft and used to help implementors identify and interoperate with deployed versions of the WebSocket protocol, such as versions developed before the publication of this RFC. "Interim" version numbers are subject to the "Expert Review" IANA registration policy [RFC5226], with the chairs of the HYBI Working Group (or, if the working group closes, the Area Directors for the IETF Applications Area) being the initial Designated Experts.

IANA has added initial values to the registry as follows.

Version Number	Reference	Status
0	+ draft-ietf-hybi-thewebsocketprotocol-00	Interim
1	+ draft-ietf-hybi-thewebsocketprotocol-01	Interim
2	+ draft-ietf-hybi-thewebsocketprotocol-02	Interim
3	+ draft-ietf-hybi-thewebsocketprotocol-03	Interim
4	+ draft-ietf-hybi-thewebsocketprotocol-04	Interim
5	+ draft-ietf-hybi-thewebsocketprotocol-05	Interim
6	+ draft-ietf-hybi-thewebsocketprotocol-06	Interim
7	+ draft-ietf-hybi-thewebsocketprotocol-07	Interim
8	+ draft-ietf-hybi-thewebsocketprotocol-08	Interim
9	+ Reserved	
10	+ Reserved	
11	+ Reserved	
12	+ Reserved	
13	+ RFC 6455	Standard

## 11.7. WebSocket Close Code Number Registry

This specification creates a new IANA registry for WebSocket Connection Close Code Numbers in accordance with the principles set out in RFC 5226 [RFC5226].

As part of this registry, IANA maintains the following information:

Status Code	The Status Code denotes a reason for a WebSocket connection closure as per <a href="#">Section 7.4</a> of this document. The status code is an integer number between 1000 and 4999 (inclusive).
Meaning	The meaning of the status code. Each status code has to have a unique meaning.
Contact	A contact for the entity reserving the status code.
Reference	The stable document requesting the status codes and defining their meaning. This is required for status codes in the range 1000-2999 and recommended for status codes in the range 3000-3999.

WebSocket Close Code Numbers are subject to different registration requirements depending on their range. Requests for status codes for use by this protocol and its subsequent versions or extensions are subject to any one of the "Standards Action", "Specification Required" (which implies "Designated Expert"), or "IESG Review" IANA registration policies and should be granted in the range 1000-2999. Requests for status codes for use by libraries, frameworks, and applications are subject to the "First Come First Served" IANA registration policy and should be granted in the range 3000-3999. The range of status codes from 4000-4999 is designated for Private Use. Requests should indicate whether they are requesting status codes for use

by the WebSocket Protocol (or a future version of the protocol), by extensions, or by libraries/frameworks/applications.

IANA has added initial values to the registry as follows.

Status Code	Meaning	Contact	Reference
1000	Normal Closure	hybi@ietf.org	RFC 6455
1001	Going Away	hybi@ietf.org	RFC 6455
1002	Protocol error	hybi@ietf.org	RFC 6455
1003	Unsupported Data	hybi@ietf.org	RFC 6455
1004	---Reserved---	hybi@ietf.org	RFC 6455
1005	No Status Rcvd	hybi@ietf.org	RFC 6455
1006	Abnormal Closure	hybi@ietf.org	RFC 6455
1007	Invalid frame payload data	hybi@ietf.org	RFC 6455
1008	Policy Violation	hybi@ietf.org	RFC 6455
1009	Message Too Big	hybi@ietf.org	RFC 6455
1010	Mandatory Ext.	hybi@ietf.org	RFC 6455
1011	Internal Server Error	hybi@ietf.org	RFC 6455
1015	TLS handshake	hybi@ietf.org	RFC 6455

## 11.8. WebSocket Opcode Registry

This specification creates a new IANA registry for WebSocket Opcodes in accordance with the principles set out in RFC 5226 [RFC5226].

As part of this registry, IANA maintains the following information:

Opcode	The opcode denotes the frame type of the WebSocket frame, as defined in Section 5.2. The opcode is an integer number between 0 and 15, inclusive.
Meaning	The meaning of the opcode value.
Reference	The specification requesting the opcode.

WebSocket Opcode numbers are subject to the "Standards Action" IANA registration policy [RFC5226].

IANA has added initial values to the registry as follows.

Opcode	Meaning	Reference
0	Continuation Frame	RFC 6455
1	Text Frame	RFC 6455
2	Binary Frame	RFC 6455
8	Connection Close Frame	RFC 6455
9	Ping Frame	RFC 6455
10	Pong Frame	RFC 6455

### 11.9. WebSocket Framing Header Bits Registry

This specification creates a new IANA registry for WebSocket Framing Header Bits in accordance with the principles set out in RFC 5226 [RFC5226]. This registry controls assignment of the bits marked RSV1, RSV2, and RSV3 in Section 5.2.

These bits are reserved for future versions or extensions of this specification.

WebSocket Framing Header Bits assignments are subject to the "Standards Action" IANA registration policy [RFC5226].

## 12. Using the WebSocket Protocol from Other Specifications

The WebSocket Protocol is intended to be used by another specification to provide a generic mechanism for dynamic author-defined content, e.g., in a specification defining a scripted API.

Such a specification first needs to *Establish a WebSocket Connection*, providing that algorithm with:

- The destination, consisting of a */host/* and a */port/*.
- A */resource name/*, which allows for multiple services to be identified at one host and port.
- A */secure/* flag, which is true if the connection is to be encrypted and false otherwise.
- An ASCII serialization of an origin [RFC6454] that is being made responsible for the connection.
- Optionally, a string identifying a protocol that is to be layered over the WebSocket connection.

The */host/*, */port/*, */resource name/*, and */secure/* flag are usually obtained from a URI using the steps to parse a WebSocket URI's components. These steps fail if the URI does not specify a WebSocket.

If at any time the connection is to be closed, then the specification needs to use the *Close the WebSocket Connection* algorithm (Section 7.1.1).

Section 7.1.4 defines when *The WebSocket Connection is Closed*.

While a connection is open, the specification will need to handle the cases when *A WebSocket Message Has Been Received* (Section 6.2).

To send some data */data/* to an open connection, the specification needs to *Send a WebSocket Message* (Section 6.1).

## 13. Acknowledgements

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